

Functional-Structural Plant Modelling with GroIMP and XL

Tutorial and Workshop at Agrocampus Ouest, Angers, 5-7 May, 2015

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More advanced XL programming (Queries etc.)

The language XL

- extension of Java
- allows also specification of L-systems and RGGs (graph grammars) in an intuitive rule notation

```
procedural blocks, like in Java: { ... }
```

```
rule-oriented blocks (RGG blocks): [...]
```

• nodes of the graph are Java objects (including geometry objects)

```
example: XL programme for the Koch curve (see part 1)
```

```
public void derivation()
[
Axiom ==> RU(90) F(10);
F(x) ==> F(x/3) RU(-60) F(x/3) RU(120) F(x/3) RU(-60) F(x/3);
]
nodes of the edges (type "successor")
graph
```

special nodes:

geometry objects

Box, Sphere, Cylinder, Cone, Frustum, Parallelogram...

access to attributes by parameter list:

Box(x, y, z) (length, width, height)

or with special functions:

Box(...).(setColor(0x007700)) (colour)

special nodes:

```
geometry objects
```

Box, Sphere, Cylinder, Cone, Frustum, Parallelogram... transformation nodes

Translate(x, y, z), Scale(cx, cy, cz), Scale(c),

Rotate(a, b, c), RU(a), RL(a), RH(a), RV(c), RG, ...

light sources

PointLight, DirectionalLight, SpotLight, AmbientLight

 rules organized in blocks [...], control of application by control structures

```
example: rules for the stochastic tree
Axiom ==> L(100) D(5) A;
A ==> F0 LMul(0.7) DMul(0.7)
if (probability(0.5))
  ( [ RU(50) A ] [ RU(-10) A ] )
  else
  ( [ RU(-50) A ] [ RU(10) A ] );
```



• parallel application of the rules

(can be modified: sequential mode can be switched on, see below)

• parallel execution of assignments possible (*deferred assignment*)

```
special assignment operator := besides the normal = quasi-parallel assignment to the variables x and y:
```

```
x := f(x, y);
y := g(x, y);
```

- Execution rules
 - Used to execute imperative statements for searched organs or update their attribute values, without changes in the (graph) structure

```
> XL syntax: ::>
ModuleName ::> { imperative code; }
Example:
l:Leaf ::> { l.photosynthesis(); }
L-system rule Execution rule
Internode(length, radius) ==> i:Internode ::> {
Internode(length+0.1, radius+0.01);
;
```

execution rule

A ::> { imperative code };

```
Test the examples sm09_e25.rgg, sm09_e16.rgg,
sm09_e17.gsz, sm09_e18.rgg
```

and concerning the access to node attributes: sm09_e26.rgg

- set-valued expressions (more precisely: producer instead of sets)
- graph queries to analyze the actual structure

```
example for a graph query:
```

binary tree, growth shall start only if there is enough distance to other \mathbf{F} objects

without the "if" condition

with the "if" condition





query syntax:

a query is enclosed by (* *)

The elements are given in their expected order, e.g.: (* **A A B** *) searches for a subgraph which consists of a sequence of nodes of the types **A A B**, connected by successor edges.

Queries as generalized contexts:

test the examples sm09_e28.rgg, sm09_e29.rgg, sm09_e30.rgg

Examples

Find all internodes, and print them out

```
println((* Internode *));
```

Find all newly created internodes (with age 0)

```
(* i:Internode, (i[age] == 0) *)
```

Search for all internodes with diameter > 0.01

```
(* i:Internode, (i[diameter] > 0.01) *)
```

Find all pairs with distance < 1</p>

(* f:F, g:F, ((f != g) && (distance(f, g) < 1)) *)

Find all nodes B, connected to A with a branching edge

(* A +> B *)

 aggregating operators (e.g., "sum", "mean", "empty", "forall", "selectWhereMin")

can be applied to set-valued results of a query

Queries and aggregating operators

provide possibilities to connect structure and function

example: search for all leaves which are successors of node c and sum up their surface areas



aggregation operator

Aggregate methods

- Collect multiple values and return one single value as result
- Standard aggregate operations: count, sum, empty, exist, forall, first, last, max, min, mean, selectRandomly, selectWhereMin, selectWhereMax, ...
- Can be applied to set-valued results of a query

```
count((* Leaf *))
first((* l:Leaf, (l[order] == 1 && l[rank] == 1) *))
selectRandomly((* F *))
selectWhereMax((* f:F *), (f[diameter]))
```

Examples

```
Count all segments F, longer than 1
```

```
count((* f:F, (f[length] > 1) *))
```

Search for all leaves and sum up their surface areas

```
sum((* Leaf *)[area])
sum((* Leaf *).area) // alternative
```

Sum up potentional growth rate of all growing leaves

```
sum((* l:Leaf, (l.isGrowing()) *).pgr())
```

Queries in XL

test the examples
sm09_e31.rgg,
sm09_e35.rgg,
sm09_e36.rgg

for light interception / photosynthesis:

a simple model of overshadowing

using a query referring to a geometric region in space

model approach (strongly simplifying): overshadowing of an object occurs when there are further objects in an imagined cone with its apex in the object, opened into z direction (to the sky).

example:

sm09_e42.rgg

competition of three 2-dimensional model plants for light

```
module Segment(int t, int ord) extends F0;
module TBud(int t) extends F(1, 1, 1);
module LBud extends F(0.5, 0.5, 1);
Vector3d z = new Vector3d(0, 0, 1);
protected void init()
  Axiom ==> P(2) D(5) V(-0.15) [ TBud(-4) ] RU(90) M(600) RU(-90)
                                [ TBud(0) ] RU(-90) M(1200) RU(90)
                                [ TBud(-8) ];
   ]
public void run()
   Γ
   TBud(t), (t < 0) ==> TBud(t+1);
   x:TBud(t), (t >= 0 && empty( (* s:Segment, (s in cone(x, z, 45)) *) ) ==>
      L(random(80, 120)) Segment(0, 0)
      [ MRel(random(0.5, 0.9)) RU(60) LBud ]
      [ MRel(random(0.5, 0.9)) RU(-60) LBud ] TBud(t+1);
  y:LBud,
      (empty( (* s:Segment, (s in cone(y, z, 45)) *) ) ==>
      L(random(60, 90) Segment(0, 1) RV0 LBud;
   Segment(t, o), (t < 8) = Segment(t+1, o);
   Segment(t, o), (t >= 8 && o == 1) ==>> ; /* removal of the whole branch */
   ]
```

Representation of graphs in XL

- node types must be declared with "module"
- nodes can be all Java objects.
 In user-made module declarations, methods (functions) and additional variables can be introduced, like in Java
- notation for nodes in a graph:
 Node_type, optionally preceded by: label:
 Examples: A, Meristem(t), b:Bud
- notation for edges in a graph:
 - -edgetype->, <-edgetype-
- special edge types: successor edge: -successor->, > or (blank) branch edge: -branch->, +> or [refinement edge: />

Notations for special edge types

- > successor edge forward
- < successor edge backward
- --- successor edge forward or backward
- +> branch edge forward
- <+ branch edge backward
- -+- branch edge forward or backward
- /> refinement edge forward
- </ refinement edge backward
- --> arbitrary edge forward
- <-- arbitrary edge backward
- -- arbitrary edge forward or backward
- (cf. Kniemeyer 2008, p. 150 and 403)

```
user-defined edge types
```

```
const int xxx = EDGE_0; // oder EDGE_1, ..., EDGE_14
...
usage in the graph: -xxx->, <-xxx-, -xxx-</pre>
```

Notation of graphs in XL

example:



is represented in programme code as a:A [-e-> B C] [<-f- D] -g-> E [a]

(the representation is not unique!)

Derived relations

Relations between nodes connected by more than 1 edge

E.g.: Relation between nodes connected by several edges (one after the other) of the same type:



"transitive closure" of the original relation (edge)

XL notation

- 1-to-n repetitions: +
 - A (-edgetype->) + B
- 0-to-*n* repetitions: *
 - A $(-edgetype->) \star B$
- Minimal elements (stop searching once a match has been found):
 - A (-edgetype->) + : (B)
 - A (-edgetype->) * : (B)



ancestor - nearest preceding node of a certain node type
minDescendants - nearest successors of a certain node type
(nodes of other types are skipped)
descendants - all successors of a certain node type



// all modules extend the module Internode
Axiom ==>

InternodeFirst [RU(30) BranchFirst] InternodeSecond [RU(-30) BranchSecond] InternodeThird

// queries and their outputs:

(* InternodeFirst -minDescendants-> Internode *)
 BranchFirst
 InternodeSecond

(* InternodeFirst -descendants-> Internode *)

InternodeSecond InternodeThird BranchSecond BranchFirst

(* InternodeThird -ancestor-> Internode *) InternodeSecond

(* InternodeThird (-ancestor->)+ Internode *)
 InternodeSecond
 InternodeFirst

The current graph

GroIMP maintains always a graph which contains the complete current structural information. This graph is transformed by application of the rules.

Attention: Not all nodes are visible objects in the 3-D view of the structure!

- F0, F(x), Box, Sphere: Yes
- RU(30), A, B: normally not (if not derived by "extends" from visible objects)

The graph can be completely visualized in the **2-D graph** view (in GroIMP: Panels - 2D - Graph).

Load an example RGG file in GroIMP and execute some steps (do not work with a too complex structure).

Open the 2-D graph view, fix the window with the mouse in the GroIMP user interface and test different layouts (Layout - Edit).

Keep track of the changes of the graph when you apply the rules (click on "redraw")!



```
which parts of the current graph of GroIMP are visible (in the 3-d view) ?
```

all geometry nodes which can be accessed from the root (denoted ^) of the graph by exactly one path, which consists only of "successor" and "branch" edges

How to enforce that an object is visible in any case:

```
==>> ^ Object
```